

High Performance Dynamic Traffic Assignment Based on Variational Inequality

GU Yangsong, LIANG Geyu

Mentor: Dr. Cheng LIU, Dr. Kwai WONG

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Outline

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- Dynamic Traffic Assignment

2. Progress

- Dynamic Network Loading Based on ODE
- Dynamic Network Loading Based on LWR
- Variational Inequality

3. Implementation

- CVODE in SUNDIAL
- Interpolators in boost library
- Ceres Library

Which runs on GPU



Heat Map Based on Vehicle Density

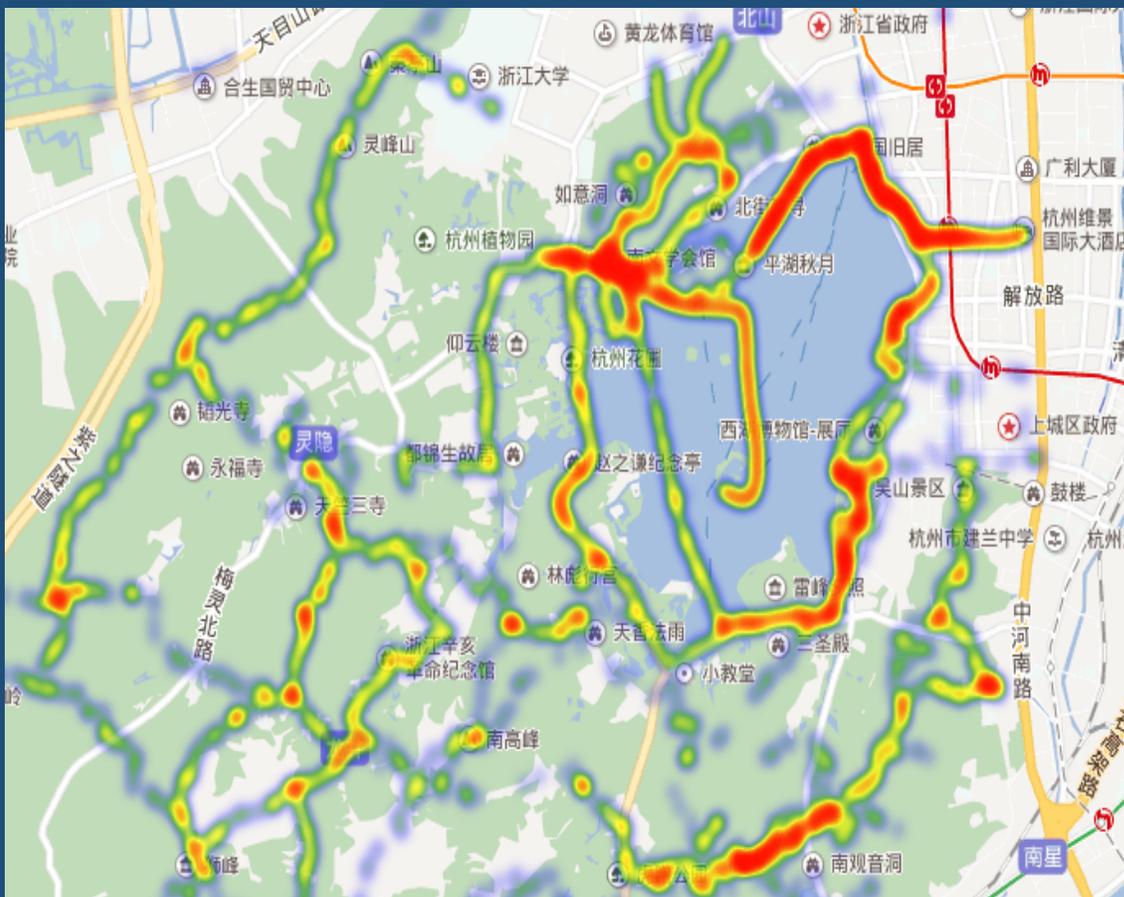
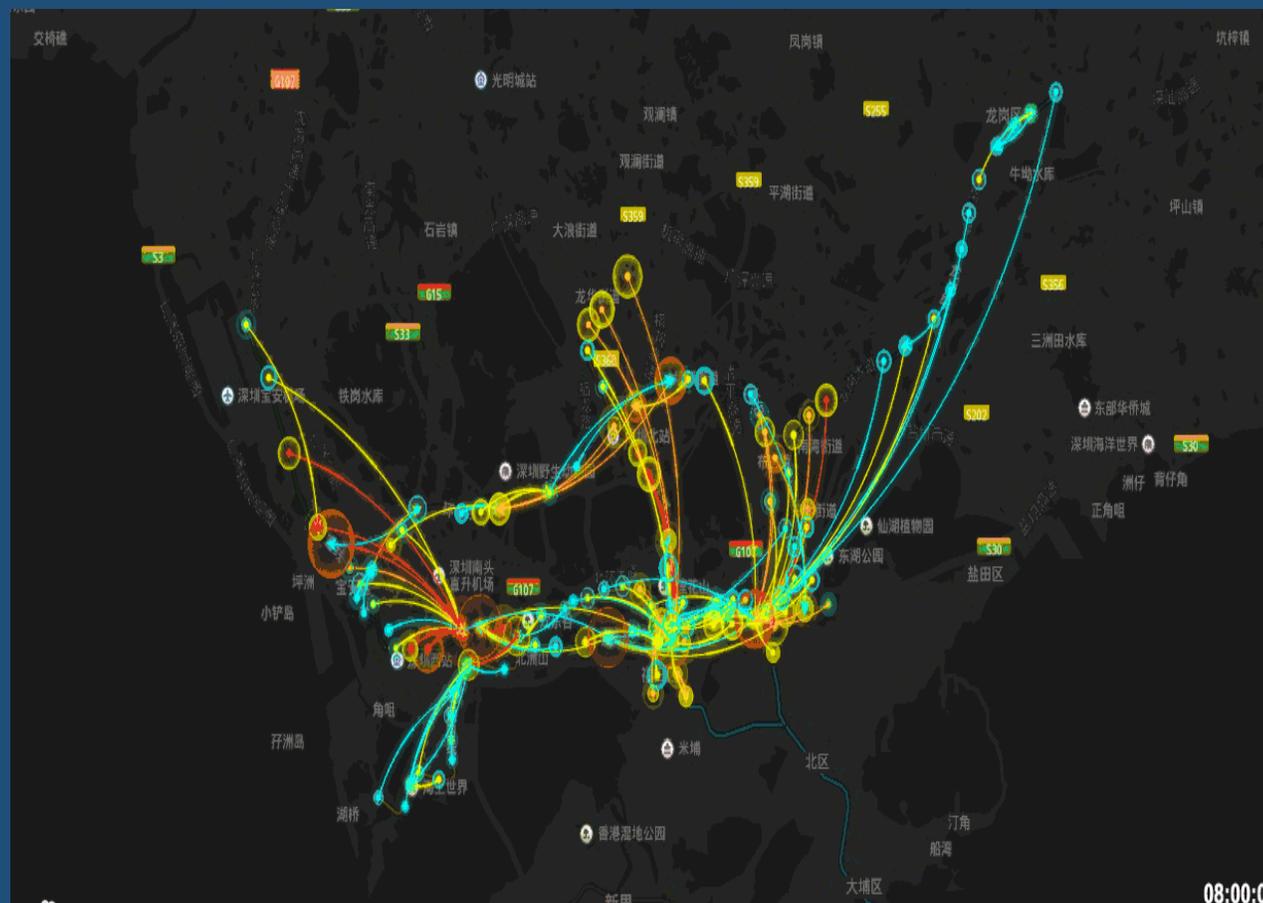
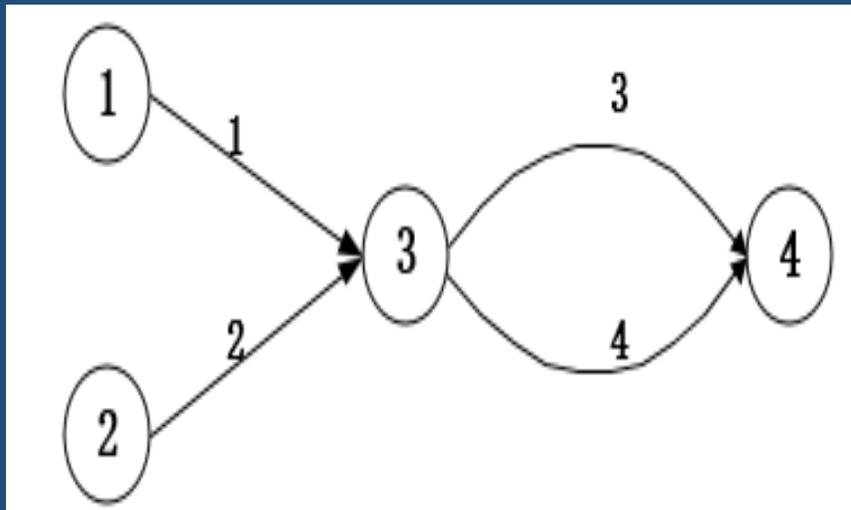


Chart of Traffic Flow



Introduction

Dynamic traffic assignment is the positive modeling of **time-varying flows** of automobiles on road network consistent with established **traffic flow theory** and **travel demand theory**.



A Simple Network

- Nodes:
- Links
- Origin-Destination Pair
- Time cost = Delay = Travel Time

Introduction

Sioux Fall Network

Departure rate function and cost function

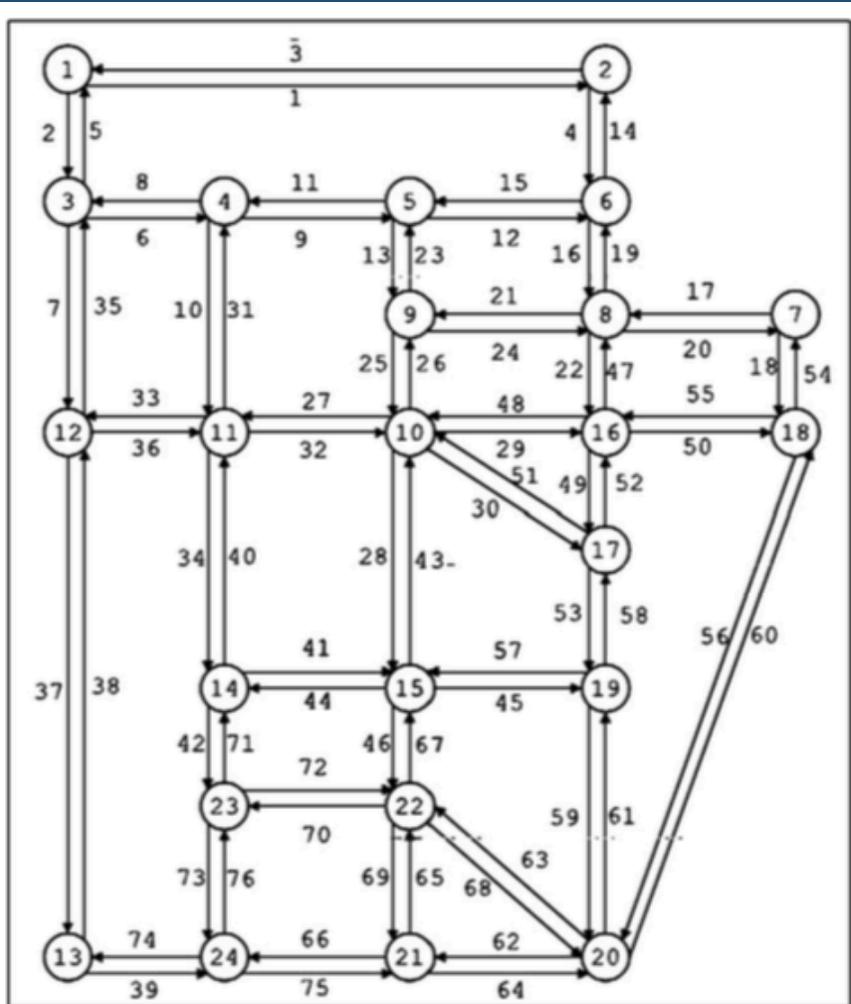
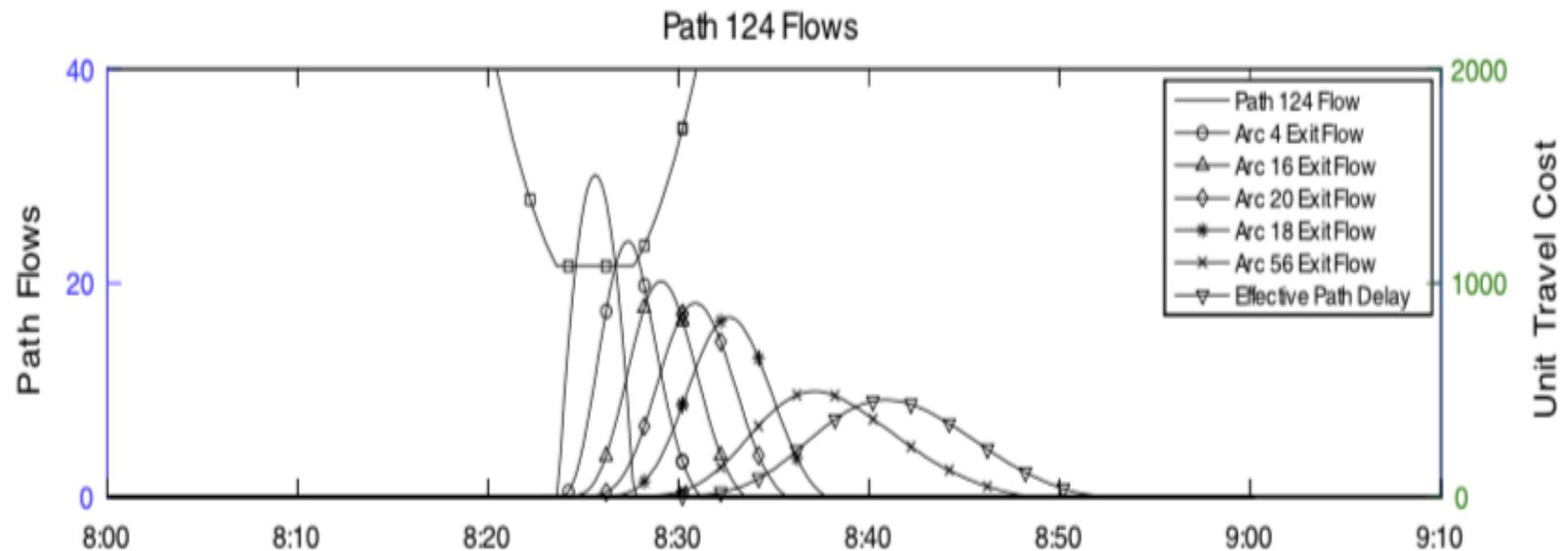


Fig. 1. Sioux Falls network.



Introduction An Optimization Problem

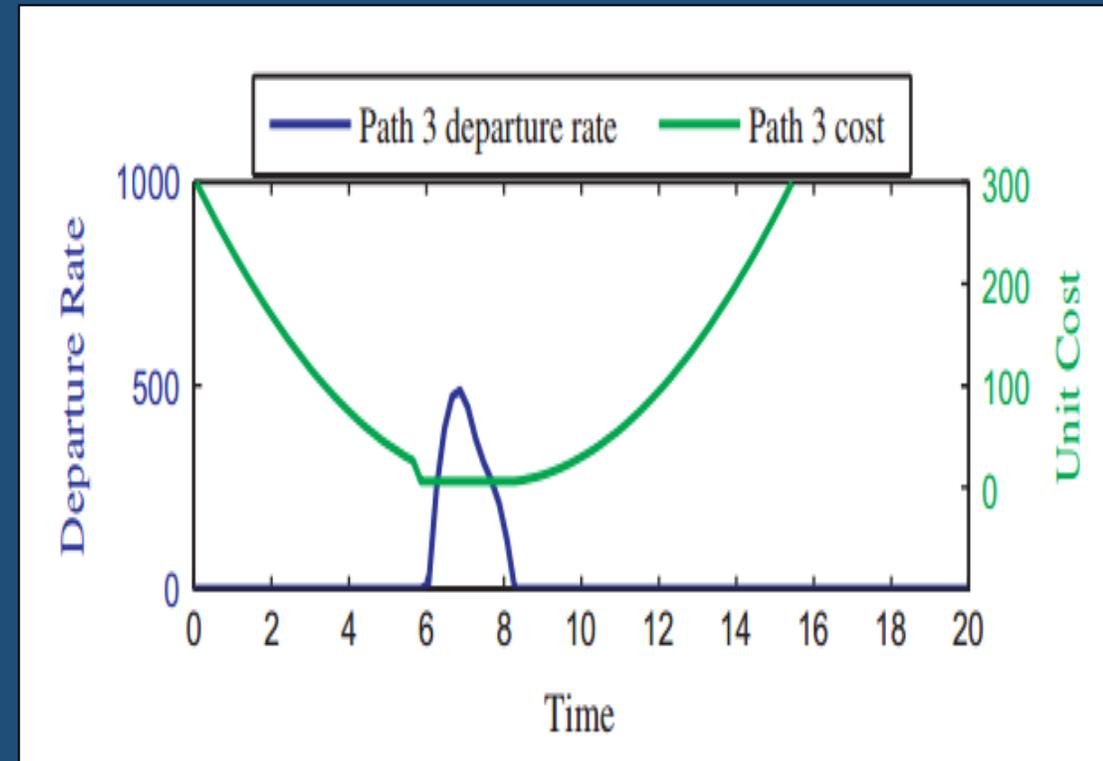
Continuous Time Dynamic User Equilibrium (DUE)

- ◆ Users choose the path with the **minimum** travel time, and the effective travel delay is **identical** for all the path and departure time of the same travel purpose.

Desired solution:

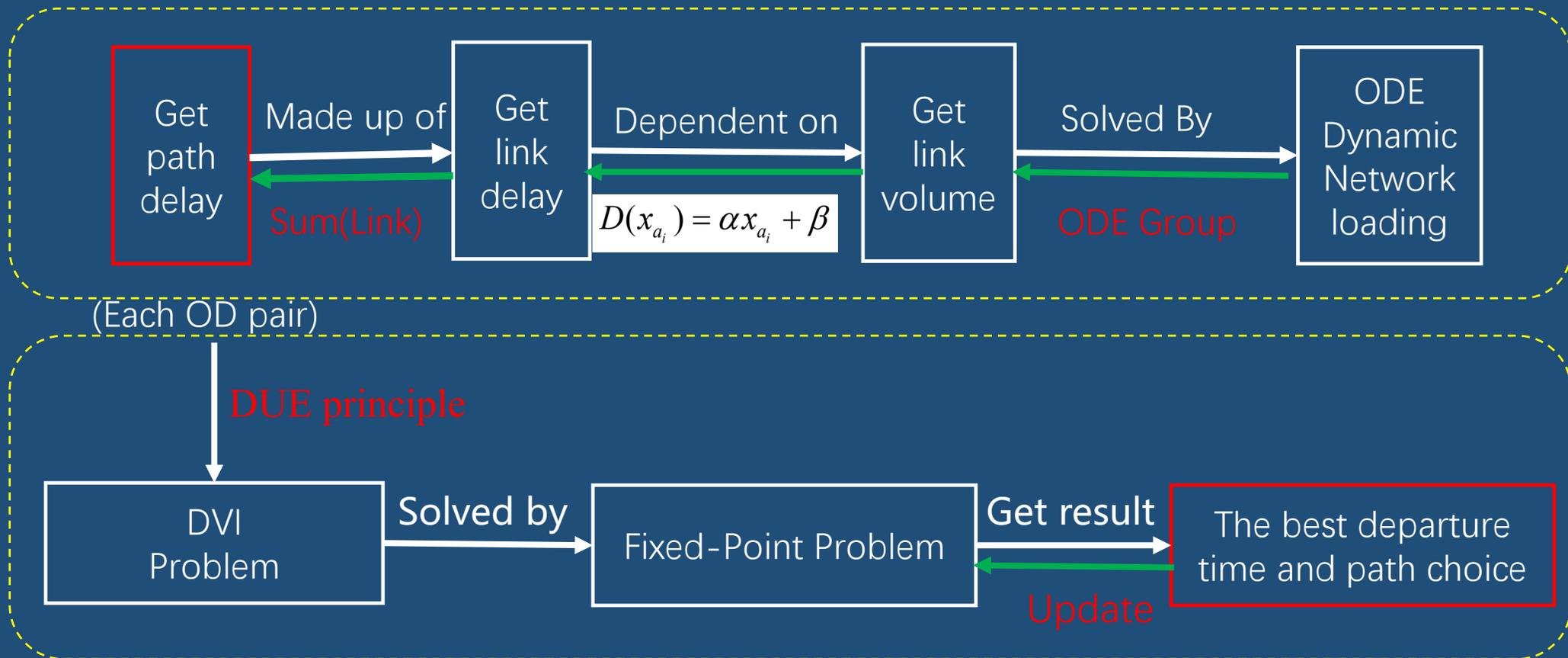
- ◆ Desired solution:
 - The **departure rate function** for each path
 - The corresponding **cost function**

Fig Departure rates and corresponding travel cost in the DUE solution



Progress

Flow chart



Progress

➤ Part 1: Dynamic Network Loading

Link state equation

$$\frac{dx_{a_1}^p(t)}{dt} = h_p^{\tau,k}(t) - g_{a_1}^p(t) \quad \forall p \in \mathcal{P}$$

$$\frac{dx_{a_i}^p(t)}{dt} = g_{a_{i-1}}^p(t) - g_{a_i}^p(t) \quad \forall p \in \mathcal{P}, i \in [2, m(p)]$$

Medium equation

$$\frac{dg_{a_i}^p(t)}{dt} = r_{a_i}^p(t) \quad \forall p \in \mathcal{P}, i \in [1, m(p)]$$

$$\frac{dr_{a_1}^p(t)}{dt} = R_{a_1}^p(x, g, r, h^{\tau,k}) \quad \forall p \in \mathcal{P}$$

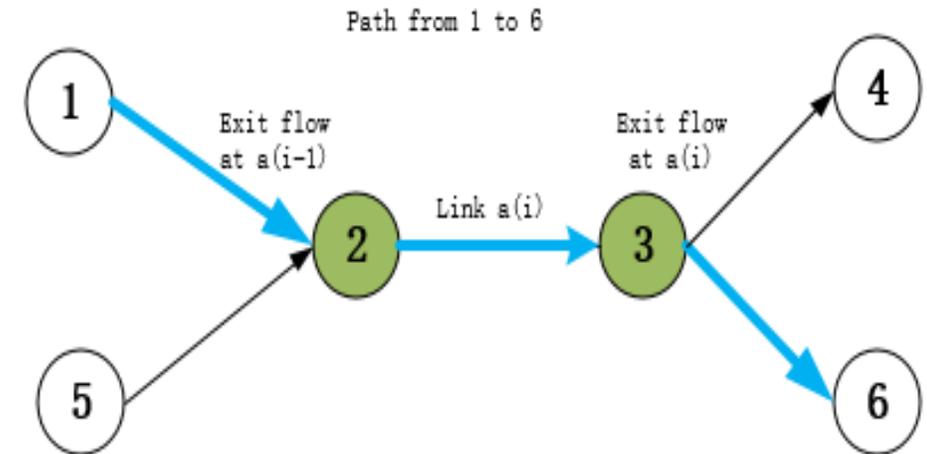
$$\frac{dr_{a_i}^p(t)}{dt} = R_{a_i}^p(x, g, r) \quad \forall p \in \mathcal{P}, i \in [2, m(p)]$$

Initial conditions

$$x_{a_i}^p((\tau - 1) \cdot \Delta) = x_{a_i}^{p,0} \quad \forall p \in \mathcal{P}, i \in [1, m(p)]$$

$$g_{a_i}^p((\tau - 1) \cdot \Delta) = 0 \quad \forall p \in \mathcal{P}, i \in [1, m(p)]$$

$$r_{a_i}^p((\tau - 1) \cdot \Delta) = 0 \quad \forall p \in \mathcal{P}, i \in [1, m(p)]$$



Flow Propagation

Progress

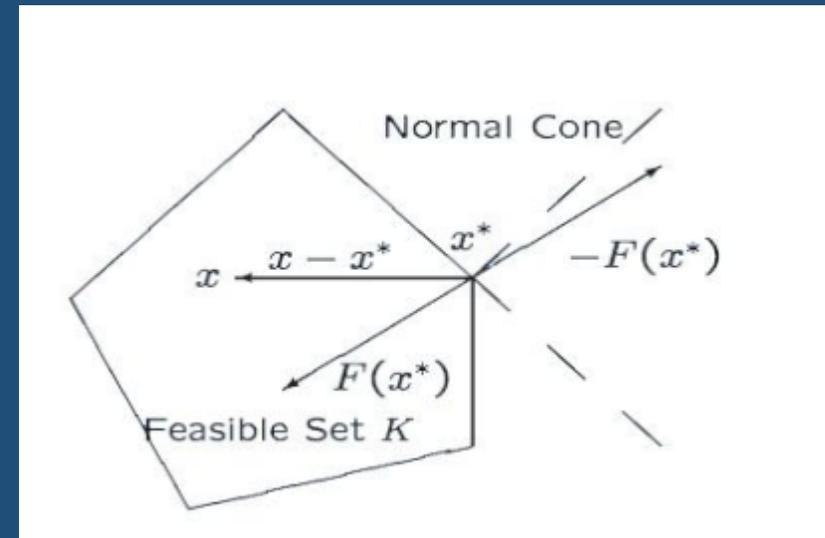
➤ Part 2: Convert DUE to DVI

$$\left. \begin{aligned}
 &\text{find } h^* \in \Lambda_0 \text{ such that} \\
 &\sum_{p \in P_{ij}} \int_{t_0}^{t_f} \Psi_p(t, h^*) (h_p - h_p^*) dt \geq 0 \quad \forall h \in \Lambda
 \end{aligned} \right\}$$

Where spatial condition is

$$\left. \begin{aligned}
 \frac{dy_{ij}}{dt} &= \sum_{p \in P_{ij}} h_p(t) & \forall (i, j) \in \mathcal{W} \\
 y_{ij}(t_0) &= 0 & \forall (i, j) \in \mathcal{W} \\
 y_{ij}(t_f) &= Q_{ij} & \forall (i, j) \in \mathcal{W}
 \end{aligned} \right\}$$

satisfy  Flow propagation constraints



Progress

➤ Part 3: Solve DVI by Fixed-Point Iteration

Equal solution

$$h^* = P_{\Lambda}[h^* - \alpha\Phi(t, h^*)]$$

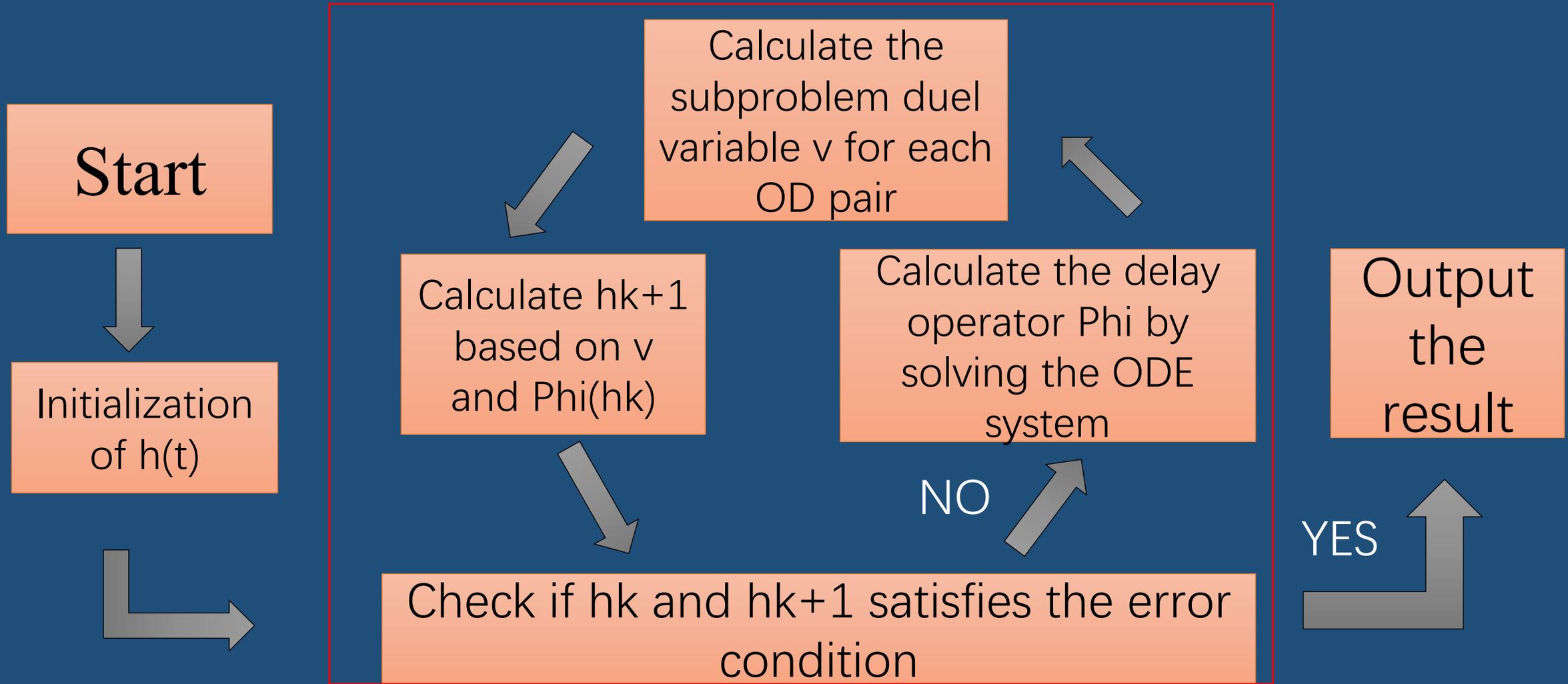
For each iteration step

$$\sum_{p \in P_{ij}} \int_{t_0}^{t_f} [h_p^k(t) - \alpha\Phi(t, h_p^k) + v_{ij}]_+ = Q_{ij}$$

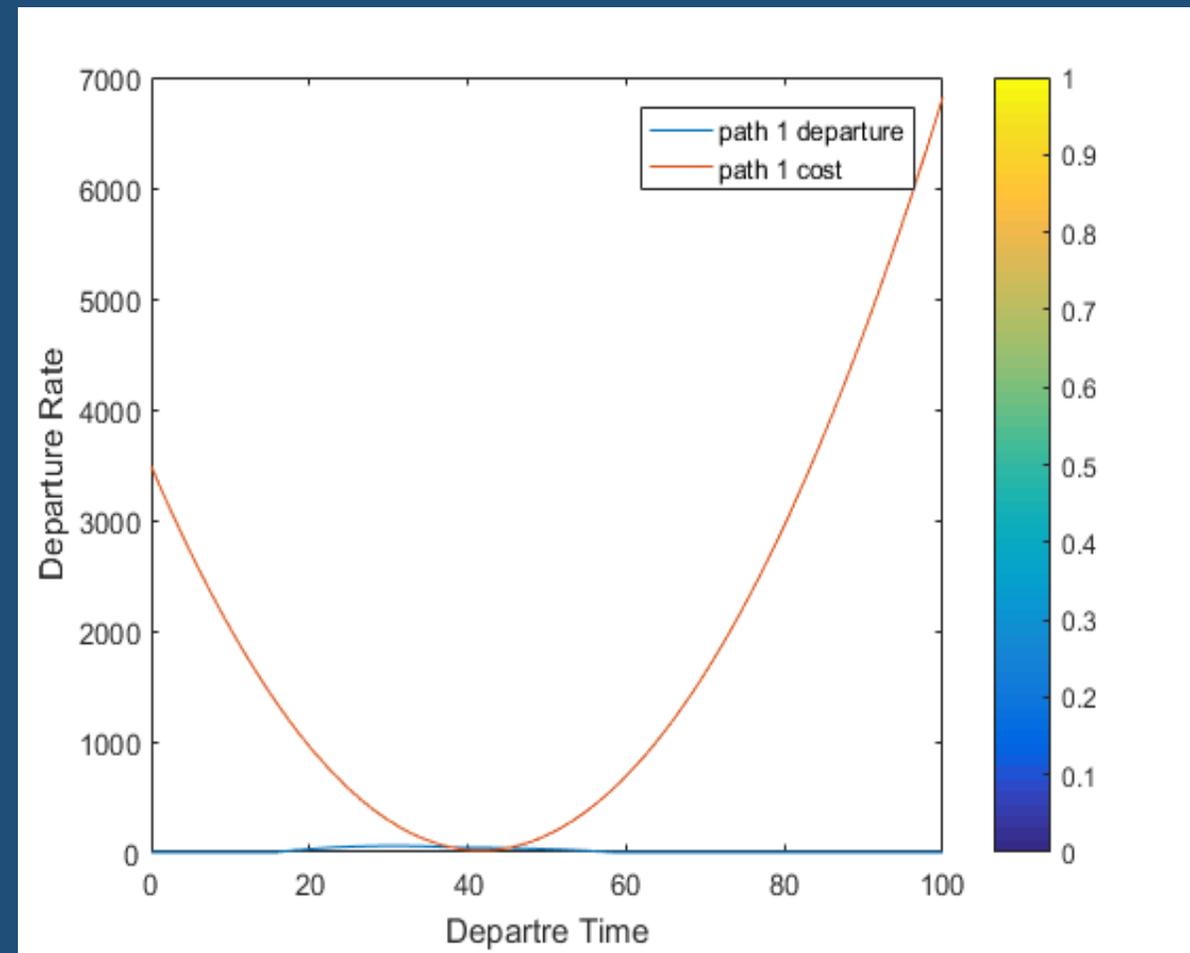
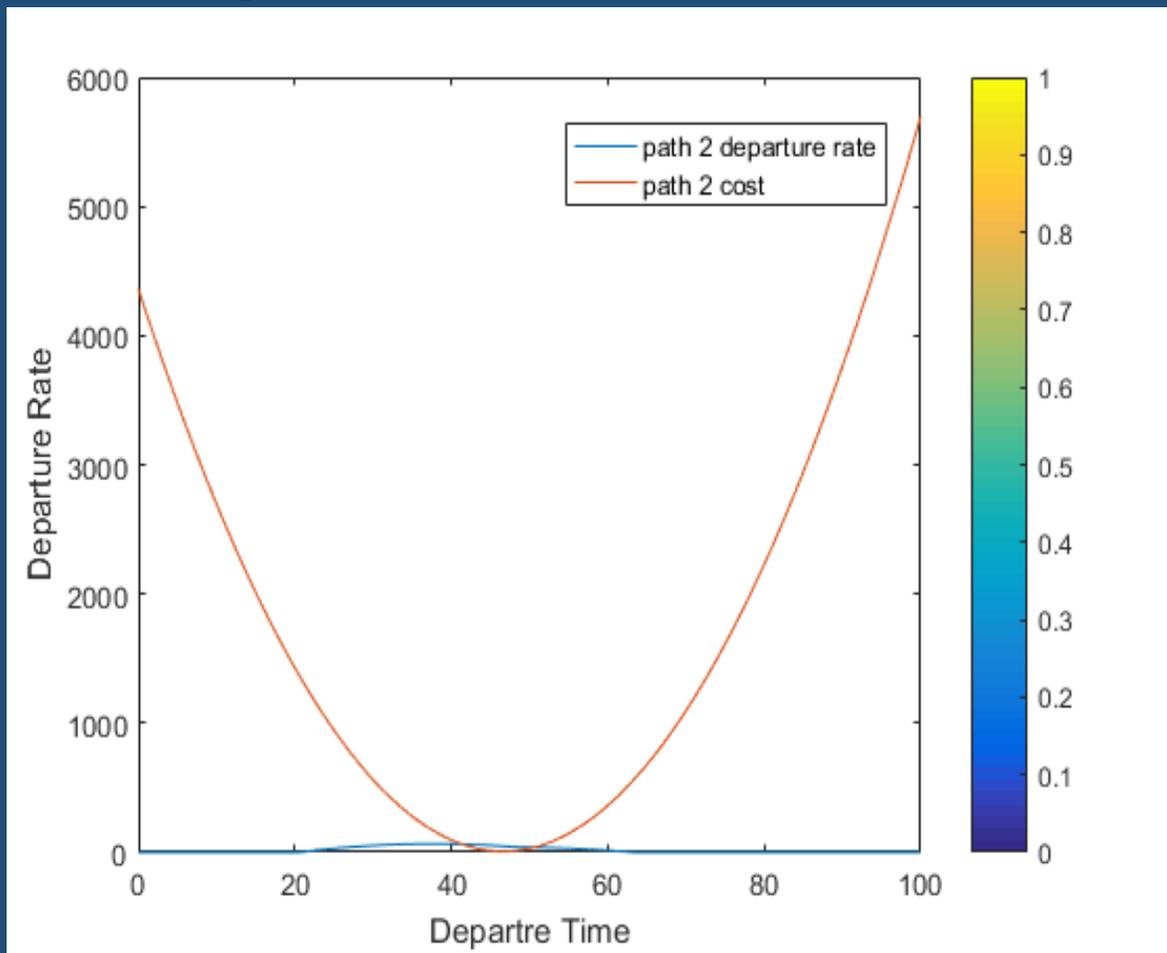
Update new departure rate

$$h_p^{k+1} = [h_p^k(t) - \alpha\Phi(t, h_p^k) + v_{ij}]_+$$

Loop of fixed point algorithm Code structure



Example



Implementation

Current Work : modify MATLAB
 convert MATLAB code to C code
 (enlarge graph scale)

ODE solver :
CVODE in SUNDIAL
odeint in boost library

Interpolator:
Interpolators in boost
library

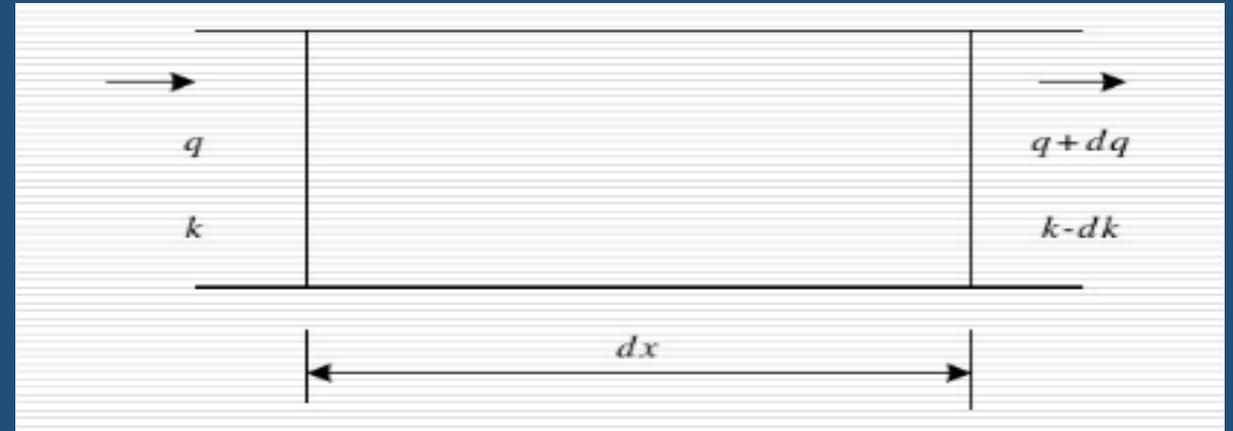
Root solver:
CERES Library

Feature Work : Dynamic Network Loading Based on LWR

Dynamic Network Loading Based on LWR

LWR model

Microcosmic Based on **PDE**

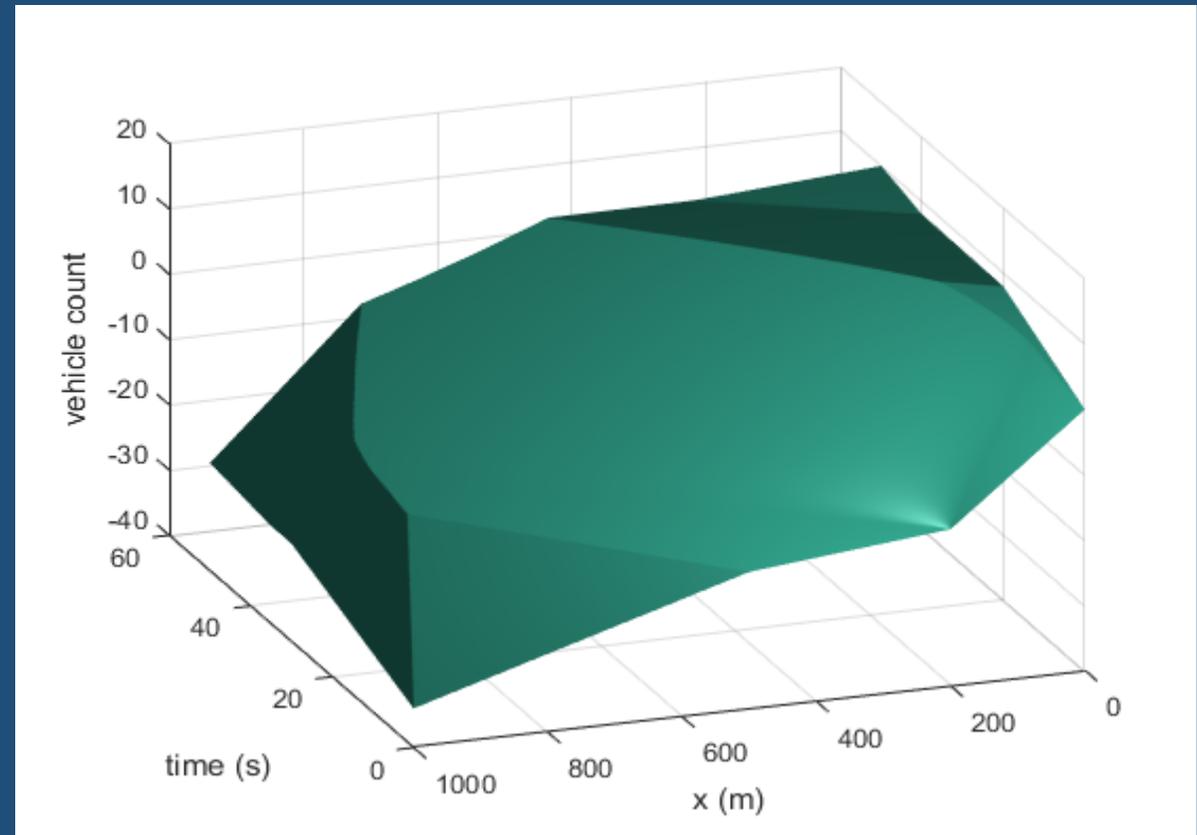
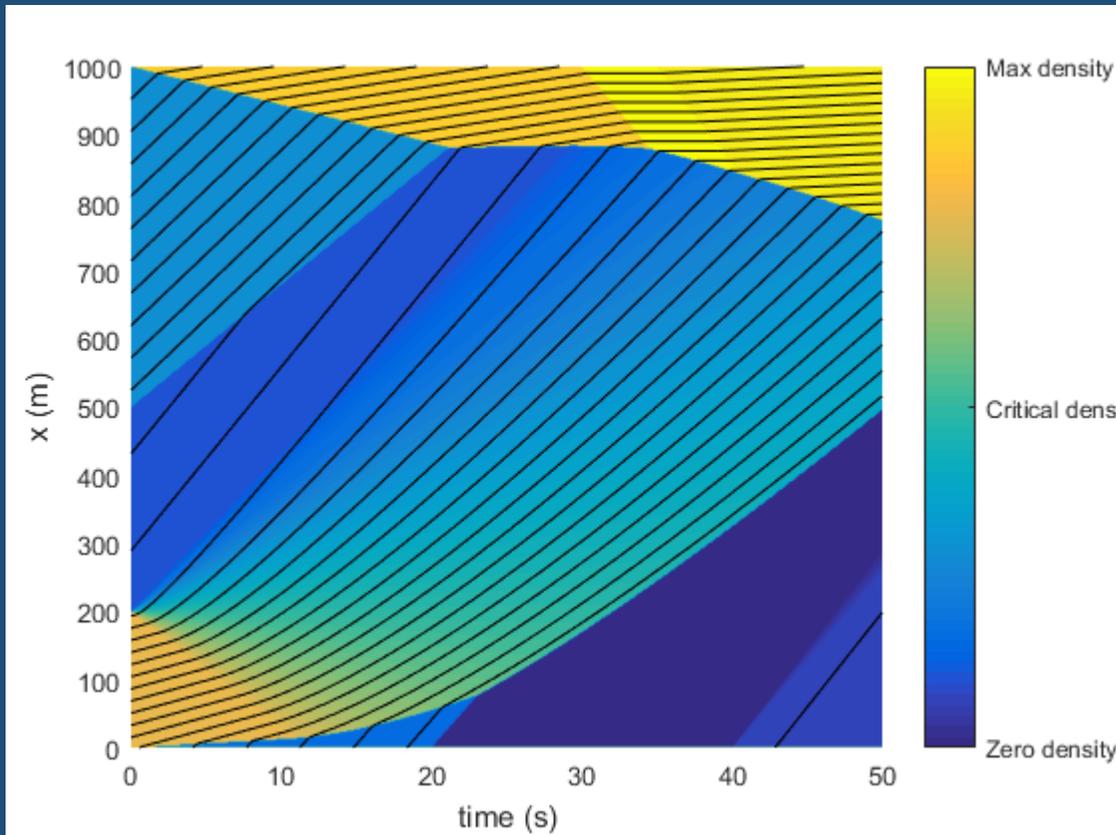


$$\begin{cases} \partial_t \rho^e(t, x) + \partial_x f^e(\rho^e(t, x)) = 0, & (t, x) \in [t_0, t_f] \times [a^e, b^e] \\ \rho^e(0, x) = 0, & x \in [a^e, b^e] \end{cases}$$

Advantages capture

1. Queues and delay
2. Density-speed relationship
3. First-in-First-out principle
4. Route information

Dynamic Network Loading Based on LWR



Link Flow Propagation

Dynamic Network Loading Based on LWR

$$Q^e(t) \doteq \sum_{p \in \mathcal{P}} Q_p^e(t), \quad q^e(t) \doteq \sum_{p \in \mathcal{P}} q_p^e(t), \quad w^e(t) \doteq \sum_{p \in \mathcal{P}} w_p^e(t)$$

$$\frac{d}{dt} Q_p^e(t) = q_p^e(t), \quad \frac{d}{dt} W^e(t) = w^e(t), \quad \forall p \in \mathcal{P}$$

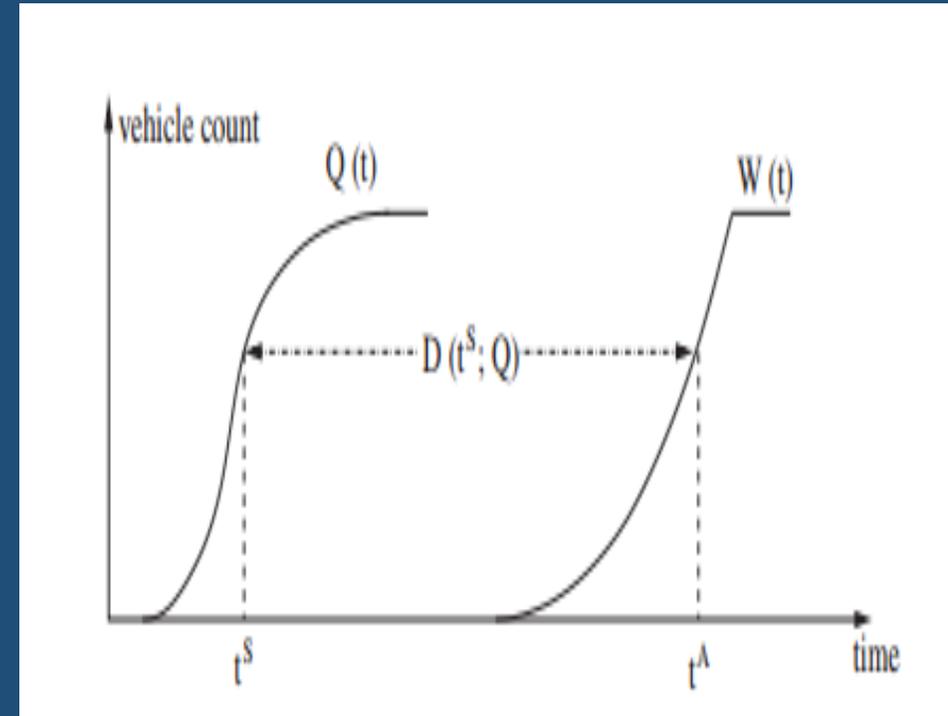
$$q_p^{e_i}(t) = w_p^{e_{i-1}}(t), \quad i \in [1, m(p)], \quad p \in \mathcal{P}$$

The solution of PDE in the form of Lax-Hopf formula

$$W^e(t) = \min_{\tau} \left\{ Q^e(\tau) + L^e \psi^e \left(\frac{t - \tau}{L^e} \right) \right\}, \quad \forall e \in \mathcal{A}$$

$$Q^e(t) = W^e(t + D(t; Q^e)), \quad \forall e \in \mathcal{A}$$

$$w_p^{e_i}(t + D(t; Q^{e_i})) = \frac{q_p^{e_i}(t)}{q^{e_i}(t)} w^{e_i}(t + D(t; Q^{e_i})), \quad i \in [1, m(p)], \quad p \in \mathcal{P}$$



Network Loading

Question

Link delay

Reference

- Friesz, T. L., Kim, T., Kwon, C., & Rigdon, M. A. (2011). Approximate network loading and dual-time-scale dynamic user equilibrium. *Transportation Research Part B: Methodological*, 45(1), 176-207. doi:10.1016/j.trb.2010.05.003
- Friesz, T. L., Han, K., Neto, P. A., Meimand, A., & Yao, T. (2013). Dynamic user equilibrium based on a hydrodynamic model. *Transportation Research Part B: Methodological*, 47, 102-126. doi:10.1016/j.trb.2012.10.001

End

Thank You